

Funcons-beta: Patterns *

The PLaNCompS Project

Patterns.cbs | PLAIN | PRETTY

OUTLINE

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Patterns

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General patterns are simple patterns or structured patterns. Matching a pattern to a value either computes an environment or fails.

Simple patterns are constructed from abstractions whose bodies depend on a given value, and whose executions either compute environments or fail.

Structured patterns are composite values whose components may include simple patterns as well as other values.

Matching a structured value to a structured pattern is similar to assigning a structured value to a structured variable, with simple pattern components matching component values analogously to simple variable components assigned component values.

Note that patterns match only values, not (empty or proper) sequences.

Meta-variables $T, T' <: \text{values}$

*Suggestions for improvement: plancomps@gmail.com.
Reports of issues: <https://github.com/plancomps/CBS-beta/issues>.

Simple patterns

Datatype `patterns ::= pattern(_ : abstractions(values \Rightarrow environments))`

`patterns` is the type of simple patterns that can match values of a particular type.

`pattern(abstraction(X))` constructs a pattern with dynamic bindings, and `pattern(closure(X))` computes a pattern with static bindings. However, there is no difference between dynamic and static bindings when the pattern is matched in the same scope where it is constructed.

Funcon `pattern-any : \Rightarrow patterns`
`~> pattern(abstraction(map()))`

`pattern-any` matches any value, computing the empty environment.

Funcon `pattern-bind(l : identifiers) : \Rightarrow patterns`
`~> pattern(abstraction(bind-value(l , given)))`

`pattern-bind(l)` matches any value, computing the environment binding l to that value.

Funcon `pattern-type(T) : \Rightarrow patterns`
`~> pattern(abstraction(if-true-else(is-in-type(given, T), map(), fail)))`

`pattern-type(T)` matches any value of type T , computing the empty environment.

Funcon `pattern-else(_ : values, _ : values) : \Rightarrow patterns`
Rule `pattern-else(P_1 : values, P_2 : values) ~>`
`pattern(abstraction(else(match(given, P_1), match(given, P_2))))`

`pattern-else(P_1, P_2)` matches all values matched by P_1 or by P_2 . If a value matches P_1 , that match gives the computed environment; if a value does not match P_1 but matches P_2 , that match gives the computed environment; otherwise the match fails.

Funcon `pattern-unite(_ : values, _ : values) : \Rightarrow patterns`
Rule `pattern-unite(P_1 : values, P_2 : values) ~>`
`pattern(abstraction(collateral(match(given, P_1), match(given, P_2))))`

`pattern-unite(P_1, P_2)` matches all values matched by both P_1 and P_2 , then uniting the computed environments, which fails if the domains of the environments overlap.

Pattern matching

Funcon `match(_ : values, _ : values) : \Rightarrow environments`

`match(V, P)` takes a (potentially structured) value V and a (potentially structured) pattern P . Provided that the structure and all components of P exactly match the structure and corresponding components of V , the environments computed by the simple pattern matches are united.

Rule $\text{match}(V : \text{values}, \text{pattern}(\text{abstraction}(X))) \rightsquigarrow \text{give}(V, X)$

Rule
$$\frac{l_2 \neq \text{"pattern"}}{\text{match}(\text{datatype-value}(l_1 : \text{identifiers}, V_1^* : \text{values}^*), \text{datatype-value}(l_2 : \text{identifiers}, V_2^* : \text{values}^*)) \rightsquigarrow \text{sequential}(\text{check-true}(\text{is-equal}(l_1, l_2)), \text{check-true}(\text{is-equal}(\text{length } V_1^*, \text{length } V_2^*)), \text{collateral}(\text{interleave-map}(\text{match}(\text{tuple-elements}(\text{given})), \text{tuple-zip}(\text{tuple}(V_1^*), \text{tuple}(V_2^*))))))}$$

Rule
$$\frac{\text{dom}(M_2) == \{ \}}{\text{match}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-equal}(\text{dom}(M_1), \{ \}), \text{map}(\text{ }, \text{fail}), \text{dom}(M_2) \neq \{ \})}$$

Rule
$$\frac{\text{some-element}(\text{dom}(M_2)) \rightsquigarrow K}{\text{match}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-in-set}(K, \text{dom}(M_1)), \text{collateral}(\text{match}(\text{map-lookup}(M_1, K), \text{map-lookup}(M_2, K)), \text{match}(\text{map-delete}(M_1, \{K\}), \text{map-delete}(M_2, \{K\}))), \text{fail})}$$

Rule
$$\frac{P : \sim(\text{datatype-values} \mid \text{maps}(-, -))}{\text{match}(V : \text{values}, P : \text{values}) \rightsquigarrow \text{if-true-else}(\text{is-equal}(V, P), \text{map}(\text{ }, \text{fail}))}$$

Funcon $\text{match-loosely}(_ : \text{values}, _ : \text{values}) : \Rightarrow \text{environments}$

$\text{match-loosely}(V, P)$ takes a (potentially structured) value V and a (potentially structured) pattern P . Provided that the structure and all components of P loosely match the structure and corresponding components of V , the environments computed by the simple pattern matches are united.

Rule `match-loosely(V : values, pattern(abstraction(X))) \rightsquigarrow give(V, X)`

Rule
$$\frac{l_2 \neq \text{"pattern"}}{\text{match-loosely}(\text{datatype-value}(l_1 : \text{identifiers}, V_1^* : \text{values}^*), \text{datatype-value}(l_2 : \text{identifiers}, V_2^* : \text{values}^*)) \rightsquigarrow \text{sequential}(\text{check-true}(\text{is-equal}(l_1, l_2)), \text{check-true}(\text{is-equal}(\text{length } V_1^*, \text{length } V_2^*)), \text{collateral}(\text{interleave-map}(\text{match-loosely}(\text{tuple-elements}(\text{given})), \text{tuple-zip}(\text{tuple}(V_1^*), \text{tuple}(V_2^*))))))}$$

Rule
$$\frac{\text{dom}(M_2) == \{ \}}{\text{match-loosely}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{map}()}$$

Rule
$$\frac{\text{dom}(M_2) \neq \{ \} \quad \text{some-element}(\text{dom}(M_2)) \rightsquigarrow K}{\text{match-loosely}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-in-set}(K, \text{dom}(M_1)), \text{collateral}(\text{match-loosely}(\text{map-lookup}(M_1, K), \text{map-lookup}(M_2, K)), \text{match-loosely}(\text{map-delete}(M_1, \{K\}), \text{map-delete}(M_2, \{K\}))), \text{fail})}$$

Rule
$$\frac{P : \sim(\text{datatype-values} \mid \text{maps}(-, -))}{\text{match-loosely}(DV : \text{values}, P : \text{values}) \rightsquigarrow \text{if-true-else}(\text{is-equal}(DV, P), \text{map}(), \text{fail})}$$

Funcon `case-match(_ : values, _ : $\Rightarrow T'$) : $\Rightarrow T'$`

`case-match(P, X)` matches P exactly to the given value. If the match succeeds, the computed bindings have scope X .

Rule `case-match(P : values, X) \rightsquigarrow scope(match(given, P), X)`

Funcon `case-match-loosely(_ : values, _ : $\Rightarrow T'$) : $\Rightarrow T'$`

`case-match(P, X)` matches P loosely to the given value. If the match succeeds, the computed bindings have scope X .

Rule `case-match-loosely(P : values, X) \rightsquigarrow scope(match-loosely(given, P), X)`

Funcon `case-variant-value(_ : identifiers) : \Rightarrow values`

`case-variant-value(I)` matches values of variant I , then giving the value contained in the variant.

Rule `case-variant-value(I : identifiers) \rightsquigarrow case-match(variant(I, pattern-any), variant-value(given))`